

Tools for proving equivalence

V.Cheval

LSV, ENS-Cachan, CNRS

School of computer science, University of Birmingham

FMATS2, Cambridge

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SECURITY PROPERTIES

Reachability properties

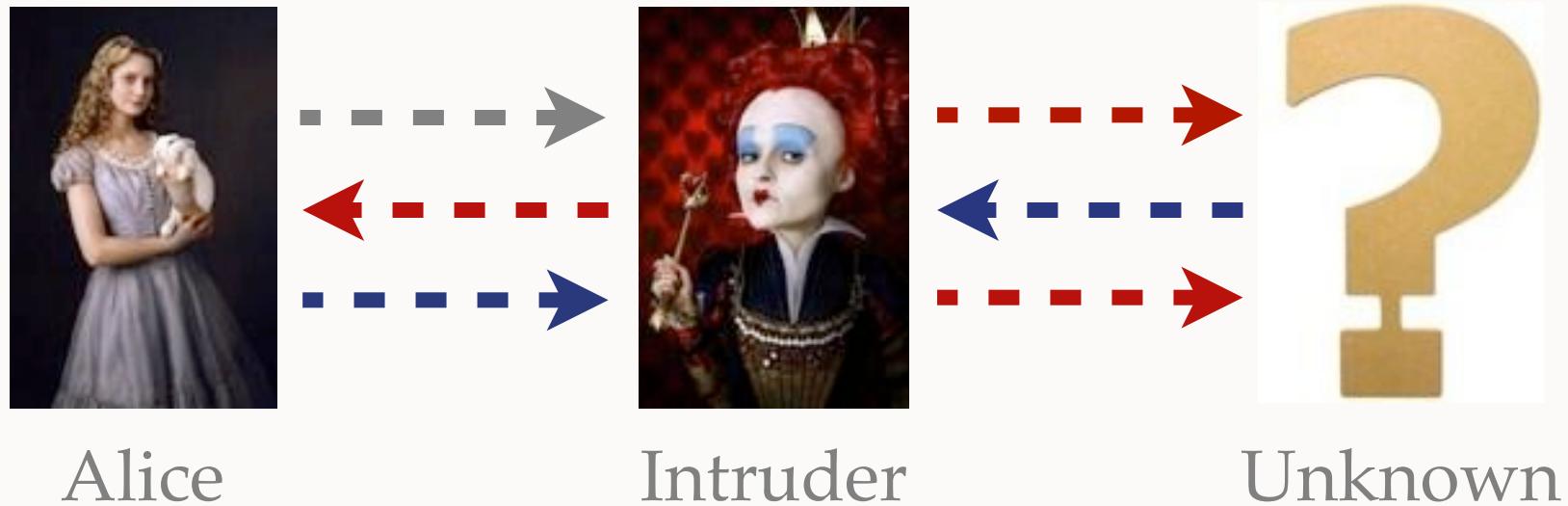
- Secrecy, Authentication, ...

Equivalence properties

- Anonymity, Privacy, Receipt-Freeness, ...

CONTEXT

- Equivalence properties : strong secret, anonymity,...

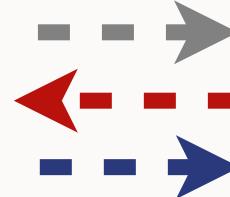


CONTEXT

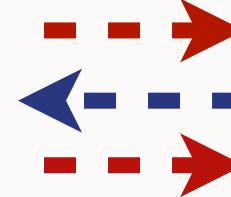
- Equivalence properties : strong secret, anonymity,...



Alice



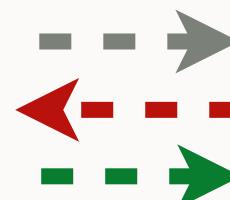
Intruder



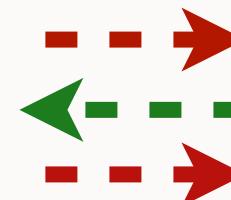
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Alice



Intruder



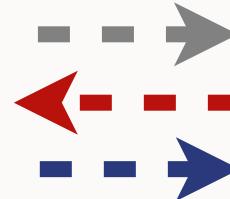
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CONTEXT

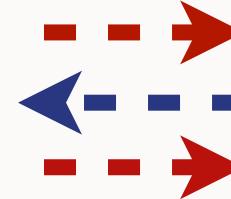
- Equivalence properties : strong secret, anonymity,...



Alice



Intruder



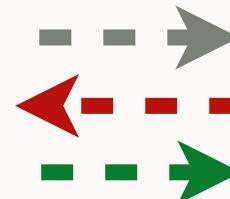
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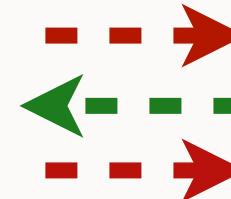
Charlene



Alice



Intruder



Unknown



Bob

TOOLS

Two problematic examples :

- e-passport protocols : M. Arapinis, T. Chothia, E. Ritter, and M. Ryan.
Analysing unlinkability and anonymity using the applied pi calculus.
- private authentication protocol : M. Abadi and C. Fournet. *Private authentication. Theoretical Computer Science.*

TOOLS

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Extension of ProVerif

- Unbounded number of sessions
- Any cryptographic primitives
- Sound
- Possible false attack
- Does not always terminate
- No more biprocess !

APTE

- Bounded number of sessions
- Fixed set of cryptographic primitives
- Sound
- Complete
- Always terminate
- Can consider length of messages

TOOLS

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Useful to prove

APTE

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- Fixed set of cryptographic primitives
- Sound
- Complete
- Always terminate
- Can consider length of messages

Useful to find attacks

APTE

- Length of messages

$$\text{enc}(x, y)$$

$$\ell_{\text{enc}}(x, y) = \alpha + \beta_1 x + \beta_2 y$$

$$\ell(\text{enc}(u, v)) = \ell_{\text{enc}}(\ell(u), \ell(v))$$

$$= \alpha + \beta_1 \ell(u) + \beta_2 \ell(v))$$